

Computability

Numeric functions

These are the elements of

$$\mathbb{N}^k \rightarrow \mathbb{N},$$

with $k \geq 0$

Notation $\text{NF}^k = \mathbb{N}^k \rightarrow \mathbb{N}$

We consider $\text{NF}^0 = \mathbb{N}$

Notation

$$\text{NF} = \bigcup_{k \in \mathbb{N}} \text{NF}^k$$

Well-known functions

$$0 \in \text{NF}^0, S \in \text{NF}^1, \text{plus} \in \text{NF}^2, \text{times} \in \text{NF}^2, \Pi_i^k \in \text{NF}^k,$$

with $\Pi_i^k(x_1, \dots, x_k) = x_i$

Primitive recursive functions

We define subsets $\text{PR}^k \subseteq \text{NF}^k$ and write

$$\text{PR} = \bigcup_k \text{PR}^k$$

the set of *primitive recursive functions* as the least set such that

Initial functions $0 \in \text{PR}^0$

$S \in \text{PR}^1$

$\Pi_i^k \in \text{PR}^k$

Composition Let $G_1, \dots, G_m \in \text{PR}^k$, $H \in \text{PR}^m$; define $F \in \text{NF}^k$ by

$$F(\vec{x}) = H(G_1(\vec{x}), \dots, G_m(\vec{x}))$$

then $F \in \text{PR}^k$

Primitive recursion Let $G \in \text{PR}^k$, $H \in \text{PR}^{k+2}$; define $F \in \text{NF}^{k+1}$ by

$$F(0, \vec{x}) = G(\vec{x})$$

$$F(y+1, \vec{x}) = H(y, F(y, \vec{x}), \vec{x})$$

then $F \in \text{PR}^{k+1}$

Primitive recursive relations

Let $R \subseteq \mathbb{N}^k$.

(i) The *characteristic function* of R is $\chi_R \in \text{NF}^k$ defined by

$$\begin{aligned}\chi_R(\vec{x}) &= 0, & \text{if } R(\vec{x}), \\ &= 1, & \text{else.}\end{aligned}$$

(ii) The relation R is *primitive recursive* if $\chi_R \in \text{PR}^k$

Then we write $R \in \text{PR}^k$.

Developing PR relations

Lemma. Let $R \in PR$. Then also $\neg R \in PR^k$.

Proof. $\chi_{\neg R}(\vec{x}) = \overline{sg}(\chi_R(\vec{x}))$. ■

Lemma *Let $R_1, R_2 \in PR$. Then also $R_1 \ \& \ R_2, R_1 \vee R_2 \in PR$.*

Proof. $\chi_{R_1 \vee R_2}(\vec{x}) = \chi_{R_1}(\vec{x}) \cdot \chi_{R_2}(\vec{x})$;

$R_1(\vec{x}) \ \& \ R_2(\vec{x}) = \neg(\neg R_1(\vec{x}) \vee \neg R_2(\vec{x}))$. ■

Developing PR functions

Lemma. Let $H_1, H_2, R \in PR^k$. Define $F \in NF^k$

$$\begin{aligned} F(\vec{x}) &= H_1(\vec{x}), && \text{if } R(\vec{x}) \\ &= H_2(\vec{x}), && \text{else.} \end{aligned}$$

Then $F \in PR^k$.

Proof. $F(\vec{x}) = H_1(\vec{x}) \cdot \chi_R(\vec{x}) + H_2(\vec{x}) \cdot \chi_{\neg R}(\vec{x})$. ■

Lemma. Let $H_1 \in PR^1, G \in PR^2, H \in PR^3$. Define $F \in NF^2$ by

$$F(x, y) = G(H_1(y), H_2(x, y, x))$$

Then $F \in PR$.

Proof. Write $\vec{x} = x, y$. Then

$$\begin{aligned} F(\vec{x}) &= G(H_1(\Pi_1^2(\vec{x})), H_2(\Pi_1^2(\vec{x}), \Pi_2^2(\vec{x}), \Pi_1^2(\vec{x}))) \\ &= G(K_1(\vec{x}), K_2(\vec{x})) \end{aligned}$$

Thus K_1, K_2 , and hence also F , are all in PR. ■

Lemma. $Z = (\lambda x.0) \in PR^1$. Proof. $Z(0) = 0; Z(x+1) = Z(x) = \Pi_2^2(x, Z(x))$. ■

More PR functions

Lemma $+, \times \in PR^2$.

Proof.

$$\begin{aligned}x + 0 &= x \\x + (y + 1) &= S(x + y) \\x \cdot 0 &= 0 \\x \cdot (y + 1) &= x \cdot y + x\end{aligned}$$

Lemma. Let $G \in PR^{k+1}$. Define

$$\begin{aligned}F(\vec{x}, y) &= \sum_{i < y} G(\vec{x}, i) \\H(\vec{x}, y) &= \prod_{i < y} G(\vec{x}, i).\end{aligned}$$

Then $F, H \in PR^{k+1}$.

Proof. Define

$$\begin{aligned}F(\vec{x}, 0) &= 0 \\F(\vec{x}, y + 1) &= F(\vec{x}, y) + G(\vec{x}, y) \\H(\vec{x}, 0) &= 1 \\H(\vec{x}, y + 1) &= H(\vec{x}, y) \cdot G(\vec{x}, y). \blacksquare\end{aligned}$$

More PR relations

Lemma. Let $R \in PR$. Define

$$S(\vec{x}, z) \Leftrightarrow \exists y < z. R(\vec{x}, y)$$

$$T(\vec{x}, z) \Leftrightarrow \forall y < z. R(\vec{x}, y)$$

Then $S, T \in PR$.

Proof. $\chi_S(\vec{x}, z) = \prod_{y < z} \chi_R(\vec{x}, y)$
 $T(\vec{x}, z) \Leftrightarrow \neg \exists y < z. \neg R(\vec{x}, y)$. ■

Lemma. Let $R \in PR^{k+1}$. Define $F \in NF^k$ by

$$\begin{aligned} F(\vec{x}, z) &= \mu y < z. R(\vec{x}, y), && \text{if } y \text{ exists} \\ &= z, && \text{else.} \end{aligned}$$

Then $F \in PR^k$.

Proof. $F(\vec{x}, 0) = 0$
 $F(\vec{x}, z + 1) = F(\vec{x}, z),$ if $\exists y < z. R(\vec{x}, y),$
 $= z,$ else, if $R(\vec{x}, z),$
 $= z + 1,$ else. ■

Packing and unpacking strings of numbers

There is a bijection $j : \mathbb{N}^2 \rightarrow \mathbb{N}$ that is PR with inverses $j_i \in \text{PR}^1$

$$j(j_1(x), j_2(x)) = x$$

[$j(x, y) = \frac{1}{2}(x + y)(x + y + 1) + x$ known to Cantor]

There is a bijection $j^m : \mathbb{N}^m \rightarrow \mathbb{N}$ that is PR with inverses $j_i^m \in \text{PR}^1$

$$j^m(j_1^m(x), \dots, j_m^m(x)) = x$$

[Iterate j]

There is a bijection $\langle - \rangle : \bigcup_{k \geq 0} \mathbb{N}^k \rightarrow \mathbb{N}$ and a maps $\lambda x j. (x)_j \in \text{PR}^2$, $\text{lh} \in \text{PR}^1$ such that

$$x = \langle \rangle \quad \Rightarrow \quad \text{lh}(x) = 0 \ \& \ (x)_i = 0$$

$$x = \langle y_0, \dots, y_{m-1} \rangle \quad \Rightarrow \quad \text{lh}(x) = m \ \& \ (x)_i = y_i, \quad \text{for } i < m$$

[$\langle y_0, \dots, y_{m-1} \rangle = j(m - 1, j^m(y_0, \dots, y_{m-1}))$]

For fixed m one has $(\lambda y_0 \dots y_{m-1}. \langle y_0, \dots, y_{m-1} \rangle) \in \text{PR}^m$

There is a map $* \in \text{PR}^2$ such that

$$\langle y_0, \dots, y_{m-1} \rangle * \langle z_0, \dots, z_{n-1} \rangle = \langle y_0, \dots, y_{m-1}, z_0, \dots, z_{n-1} \rangle$$

The Fibonacci function

Define $F(0) = 1$

$$F(1) = 1$$

$$F(n + 2) = F(n) + F(n + 1)$$

Then $F \in \text{PR}^1$.

Proof. Define $H(x) = \langle F(x), F(x + 1) \rangle$. Then

$$H(0) = \langle 1, 1 \rangle$$

$$H(x + 1) = \langle (H(x))_1, (H(x))_0 + (H(x))_1 \rangle$$

is PR. Hence $F(x) = (H(x))_0$ is PR. ■

Course of value recursion

Given $F \in \text{NF}^{k+1}$ define

$$\bar{F}(\vec{x}, n) = \langle F(\vec{x}, 0), \dots, F(\vec{x}, n-1) \rangle$$

Given $H \in \text{PR}^{k+2}$. Define

$$F(\vec{x}, n) = H(\bar{F}(\vec{x}, n), \vec{x}, n)$$

Then $F \in \text{PR}$.

Proof. We show first that $\bar{F} \in \text{PR}^{k+1}$. Indeed

$$\begin{aligned}\bar{F}(\vec{x}, 0) &= \langle \rangle \\ \bar{F}(\vec{x}, n+1) &= \bar{F}(\vec{x}, n) * \langle F(\vec{x}, n) \rangle \\ &= \bar{F}(\vec{x}, n) * \langle H(\bar{F}(\vec{x}, n), \vec{x}, n) \rangle\end{aligned}$$

Then

$$F(\vec{x}, n) = (\bar{F}(\vec{x}, n+1))_n.$$

Therefore $F \in \text{PR}$. ■

Partial recursive functions

A *partial function* $\psi : \mathbb{N}^k \rightarrow \mathbb{N}$ is a relation $\psi \subseteq \mathbb{N}^k \times \mathbb{N}$ such that

$$(x, y), (x, y') \in \psi \Rightarrow y = y'$$

If $(x, y) \in \psi$, then we write $\psi(x) \downarrow$ ($\psi(x)$ *is defined*) and $\psi(x) \simeq y$.

If for no y one has $(x, y) \in \psi$, then we write $\psi(x) \uparrow$ ($\psi(x)$ *is undefined*)

Initial functions $0 \in \mathcal{PR}^0$

$$S \in \mathcal{PR}^1$$

$$\Pi_i^k \in \mathcal{PR}^k$$

Composition Let $\psi_1, \dots, \psi_m \in \mathcal{PR}^k$, $\chi \in \mathcal{PR}^m$; define $\varphi \in \mathcal{NF}^k$ by

$$\varphi(\vec{x}) \simeq \chi(\psi_1(\vec{x}), \dots, \psi_m(\vec{x}))$$

then $\varphi \in \mathcal{PR}^k$

Primitive recursion Let $\psi \in \mathcal{PR}^k$, $\chi \in \mathcal{PR}^{k+2}$; define $\varphi \in \mathcal{NF}^{k+1}$ by

$$\varphi(0, \vec{x}) \simeq \psi(\vec{x})$$

$$\varphi(y+1, \vec{x}) \simeq \chi(y, \varphi(y, \vec{x}), \vec{x})$$

then $\varphi \in \mathcal{PR}^{k+1}$

Minimalization Let $\chi \in \mathcal{PR}^{k+1}$ and define $\varphi(\vec{x}) \simeq \mu y. [\chi(\vec{x}, y) \simeq 0]$

then $\varphi \in \mathcal{PR}^k$

Small print about equality

In an equation

$$t \simeq s$$

with expressions that are partially defined, the meaning is

$$[t \downarrow \Rightarrow [s \downarrow \ \& \ t = s]] \ \& \ [s \downarrow \Rightarrow [t \downarrow \ \& \ t = s]]$$

We understand that

$$\begin{aligned} \chi(\psi_1(\vec{x}), \dots, \psi_k(\vec{x})) \downarrow &\Leftrightarrow \exists y_1, \dots, y_k. \\ &[\psi_1(\vec{x}) = y_1 \ \& \ \dots \ \& \ \psi_k(\vec{x}) = y_k \\ &\ \& \ \chi(y_1, \dots, y_k) \downarrow] \end{aligned}$$

$$\begin{aligned} \mu y. [\chi(\vec{x}, y) \simeq 0] \downarrow &\Leftrightarrow \exists y. [\chi(\vec{x}, y) \simeq 0 \ \& \\ &\forall z < y. \chi(\vec{x}, z) \downarrow \ \& \ \chi(\vec{x}, z) \neq 0] \end{aligned}$$