

# From Mind to Turing to Mind

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Many *numerical problems* can be answered by **computing**

“What is the area of a circle with radius 4 yd?”

Answer:  $4^2\pi \text{ yd}^2 = 50.2654824 \text{ yd}^2$

Also many *qualitative problems* may be answered by **computing**

“Are points  $A = (x_1, y_1)$ ,  $B = (x_2, y_2)$  and  $C = (x_3, y_3)$  in  $\mathbb{R}^2$  collinear  
i.e. do they lie on a straight line?”

Answer: if and only if  $(x_1 - x_3)(y_2 - y_3) = (x_2 - x_3)(y_1 - y_3)$

Also for *daily life actions* or in the *performing arts* and *sports* one needs **computations**

Musicians can train themselves to be accurate in the sub-millisecond range

Hence one needs **fast** and **accurate computations**

Bartok: Sonata, Ivry Gitlis violin

Ravel: Ondine from Gaspard de la Nuit, John Kane piano

Leibniz (1646-1716) conjectured:

All properly stated problems can be answered by computing (calculemus!)

He wanted to construct:

a universal language  $L$  for stating problems precisely

a machine  $M$  answering all these problems by computing

The first question Leibniz wanted to ask to such a machine is said to be  
“Does God exist?”

Quite daring around 1700 to ask this question to a machine!

Science Fiction Answer, F. Brown [1954]: “Now there is a God!”

If we restrict ourselves to mathematical problems, then there is such an  $L$

If we restrict ourselves to *numerical problems* then there is such an  $M$

A computer with a software package like [Mathematica](#) or [Maple](#)

However, Turing [1936]: *for qualitative mathematical problems  $M$  is impossible*

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However, Turing [1936]: *for qualitative mathematical problems  $M$  is impossible*

E.g. “Are there twin primes (differing by 2 like {11, 13}) greater than googolplex  $10^{10^{100}}$ ?”

is an example of such a qualitative problem

Turing [1936]: *For qualitative mathematical problems  $M$  is impossible*

Turing's negative answer about the qualitative problems

made possible the *positive* answer for the numerical problems

There is also a positive answer for a *subclass* of the qualitative problems

(symbolic computing, e.g. “Are points  $A, B, C$  collinear?”)

So we have

*numerical, symbolic, and qualitative mathematical* problems

the latter involving quantifiers “for all” ( $\forall$ ), “there exists” ( $\exists$ )

These have to do with mathematical infinity

Still we can answer some problems about infinity, by *proving*

Euclid: There are infinitely many primes: 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, ...

Hilbert believed that eventually everything can be known (by proving or computing)

“Instead of the sad ‘We don’t know’, my attitude is: ‘We *have* to know, we *will* know’!”

How did Turing prove his negative result?

How could it have so much positive spin off?

Turing did the following

- ! Gave a well-motivated analysis of computability via *Turing Machines* (TM)  
an idealized class of machines (with infinite memory)
- !! Constructed a *universal Turing Machine* UTM that can simulate any TM  
via software (*programs*)
- ! Formulated the *halting problem* (HP) that cannot be decided by any TM  
argument like liar paradox
- ! Concluded *qualitative mathematical problems* cannot be decided by any TM  
as the HP is one of them

Therefore Leibniz's ideal cannot be fulfilled for mathematical problems  
only be approximated, leave alone for philosophical problems

A *process* reacts to input with actions (in principle for always)

Think of

an animal that looks for food, a mate and tries to avoid being eaten

a machine that controls a non-stop factory

the operating system of a computer

A *computation* can be seen as a terminating process

Think of

an animal that has to decide whether to go into fight or flight

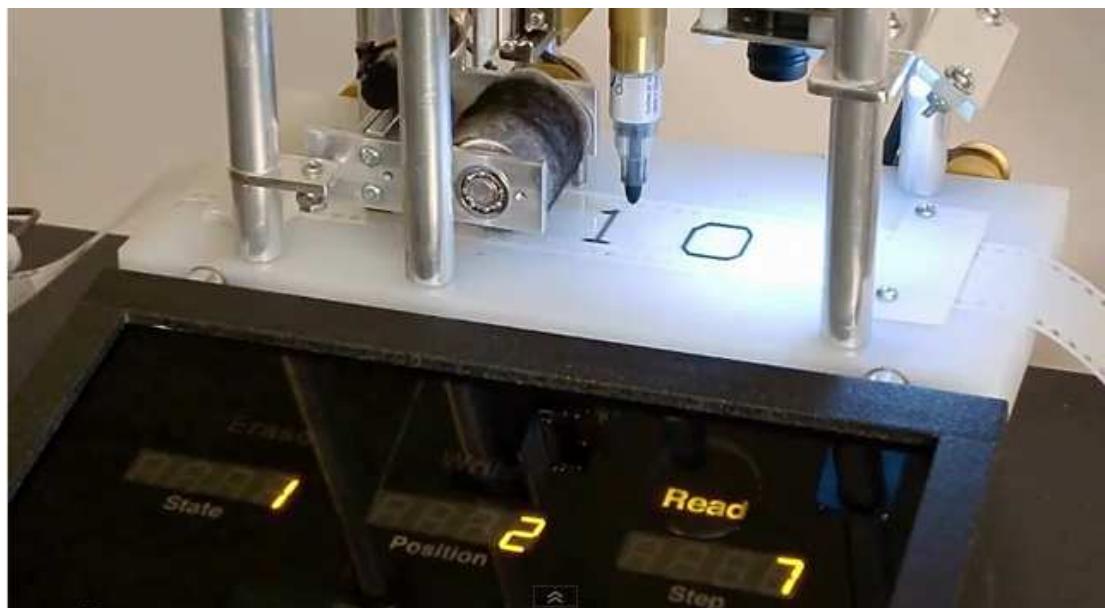
computing  $2^{127} - 1 = 170141183460469231731687303715884105727$  (a prime)

Both

computations (as terminating processes) and

proper processes (non-terminating ones)

are interesting



[www.youtube.com/watch?v=E3keLeMwfHY](https://www.youtube.com/watch?v=E3keLeMwfHY)

Based on introspection: computing is done in *discrete* steps

A computational process goes from input to output (action)

Let  $M$  be a TM. It has

a finite set  $I$  of input

a finite set  $A$  of actions

The machine  $M$  transforms an input ( $i$  in  $I$ ) into an action ( $a$  in  $A$ )

$$i \xrightarrow{M} a$$

As the same input may give rise at different times to different actions

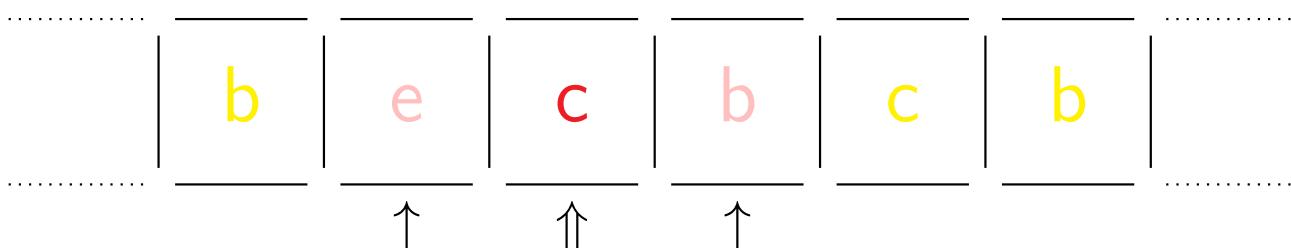
Turing introduced a set  $S$  of *states* indicating how to react to an input

$$(i, s) \xrightarrow{M} (a, s')$$

The machine, given input and state, may choose action and (a new) state

$M$  is given as a finite table of transitions  $(i, s; a, s')$

$M$  and has a two-sided infinite memory tape  
 with cells containing nothing (a blank) or a symbol  $i \in I$   
 and also has a Read/Write device ('head') placed on one of the cells  
 An (*instantaneous*) *configuration* (at a given moment)  
 is the information on the tape & the position of the head

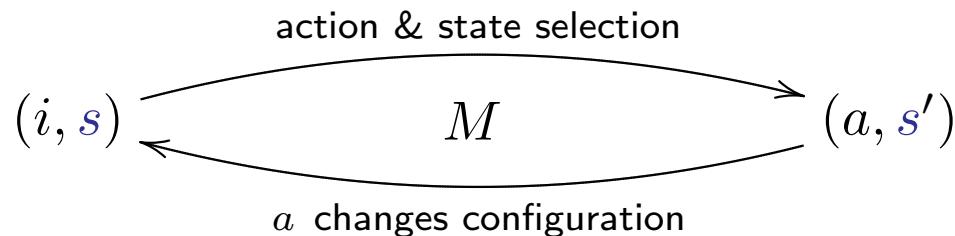


$\uparrow\uparrow$	position of read/write head	Actions
$\uparrow$	potential next position of head	$L$ $\uparrow$ goes left
<b>red</b>	letter	$R$ $\uparrow$ goes right
<b>pink</b>	letters	$W(c')$ overwrite $c$ by $c'$
<b>yellow</b>	letters	

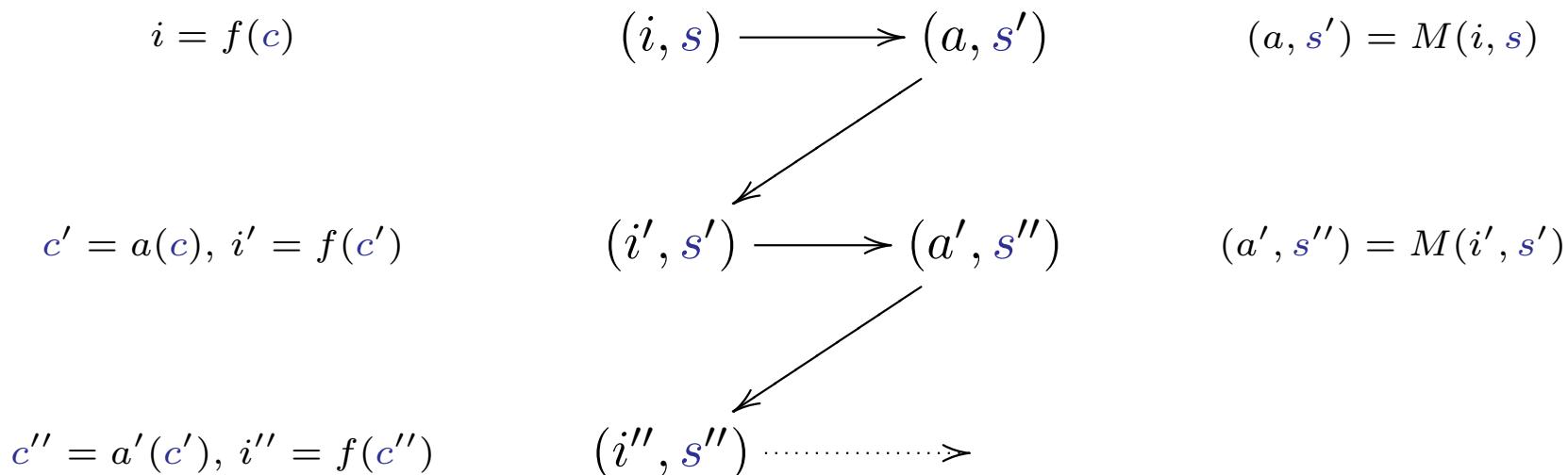
(terminology suggestive for sequel)

Actions modify configuration

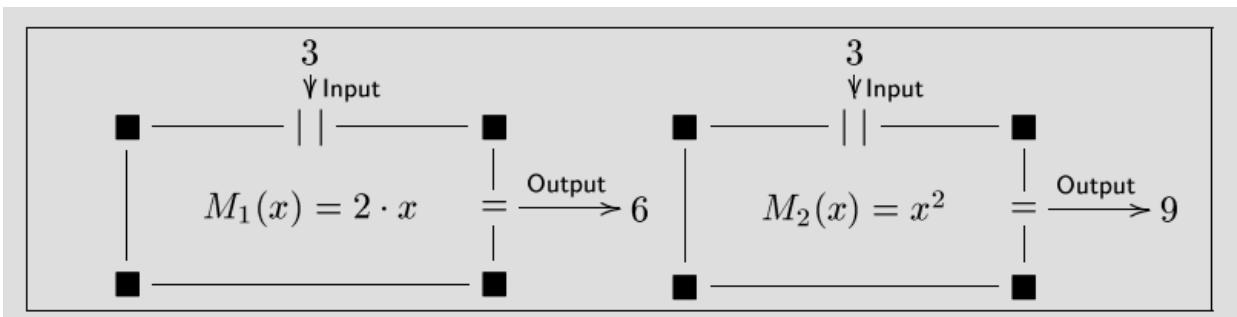
The two phases of  $M$  combine as follows



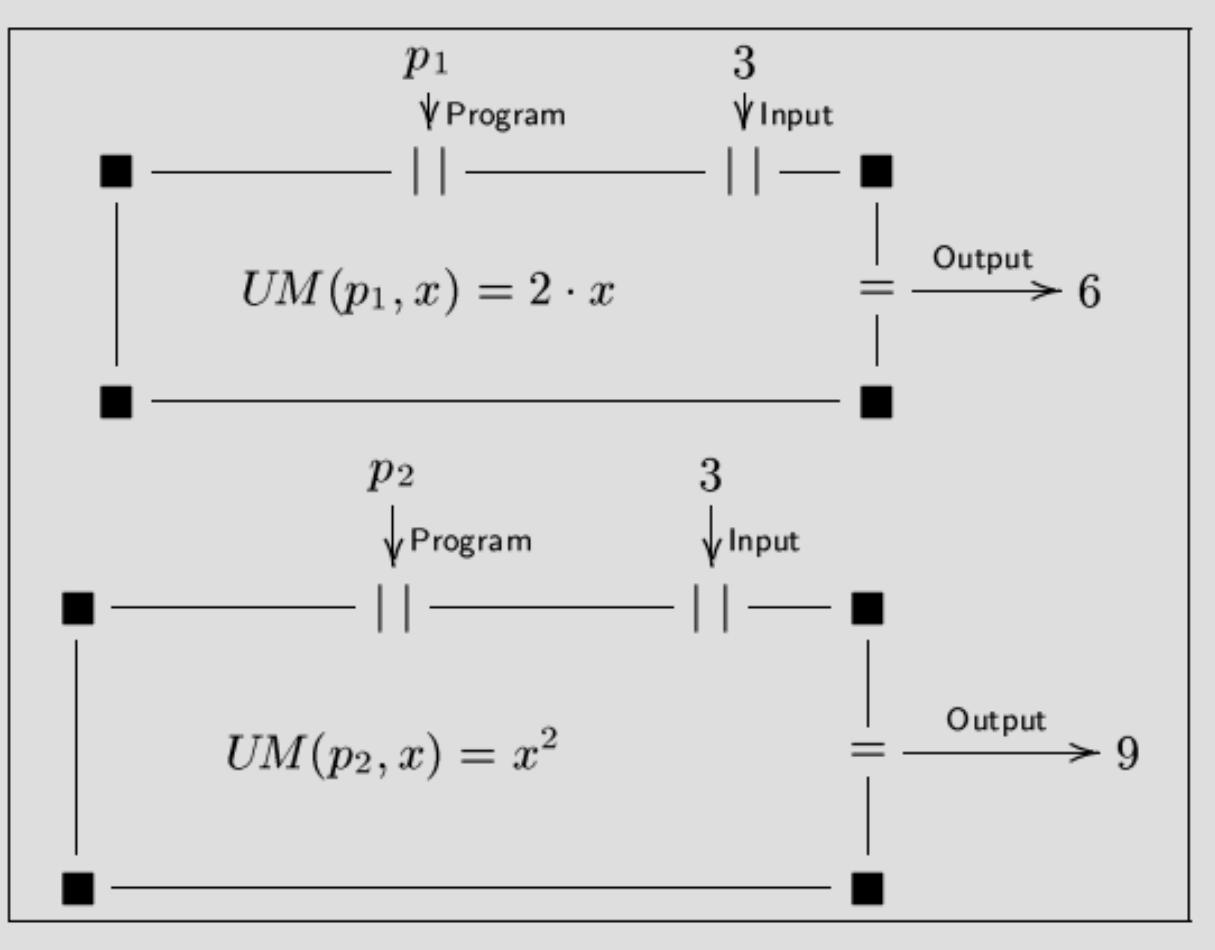
giving a ‘scenario’ as follows:  $c$  initial configuration,  $s$  initial state,  $f$  ‘focus of’



$M$  halts if no more transition is possible. Then *input*: initial  $c$ ; *output*: final  $c$



Two ad hoc machines  $M_1, M_2$



$UM$  simulating  $M_1, M_2$

Write  $M(n)\downarrow$  (resp.  $(M(n)\uparrow)$ ) if machine  $M$  with input  $n$  does (doesn't) halt

Theorem. The problem  $UM(p, n)\downarrow$  is non-computable (*undecidable*)

Proof. Suppose  $UM(p, n)\downarrow$  is computable: then there is a machine  $H$  such that

$$\begin{aligned} UM(p, n)\downarrow &\iff H(p, n) = 1 \\ UM(p, n)\uparrow &\iff H(p, n) = 0 \end{aligned}$$

Modify  $H$  into  $H'$  making  $H'(p, n)$  run forever if  $H(p, n) = 1$ . Then

$$\begin{aligned} UM(p, n)\downarrow &\iff H'(p, n)\uparrow \\ UM(p, n)\uparrow &\iff H'(p, n)\downarrow \end{aligned}$$

Now define  $D(n) = H'(n, n)$ . It has a program  $p_D$  as  $UM$  is universal:

$$D(n) = UM(p_D, n) \text{ for all } n$$

We get a contradiction

$$D(p_D)\downarrow \iff UM(p_D, p_D)\downarrow \iff H'(p_D, p_D)\uparrow \iff D(p_D)\uparrow \blacksquare$$

$UM(p, n) \downarrow \iff$  there exists a number  $k$  such that after  $k$  cycles  $UM(p, n)$  halts

The 'quantifiers'

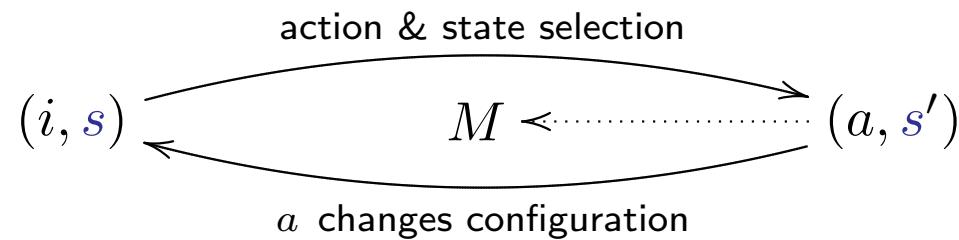
exists ( $\exists$ )

for all ( $\forall$ )

ranging over the infinite set  $\mathbb{N}$  usually cause undecidability

As any  $M$  can be written as  $M(n) = UM(p, n)$

actions may be considered also to effect the program  $p$ , hence  $M$  itself



## Memory

Rather than the tape with slow access one can use random access memory

Content addressing gives ‘associative memory’

Memory can be also divided into

core              immediately available

disc-cache      quickly available

disc/flash      taking more time

## Input/Output

In order to react to the environment some memory cells can be reserved for

values from [sensors](#)

values directing [actuators](#)

In this way a factory can be controlled by a computer  
having access to information from thermometers, barometers  
able to turn on/off a heater and exhaust valves

Nature has evolved a chemo-electrical computational model

Synapse              Neural Net

These are not programmed using intention but by trial and error

They run in parallel, can be remarkably efficient

Have the same computational power as the Turing Machine model

In human cognition there are zillion possible objects

To program how one reacts to them is undoable

But the following hybrid model [1], [2] has potential

One can see human cognition as a TM with huge  $I, A$  sets

in which the transitions are programmed by a NN

Moreover human memory is associative

The configuration can be defined by as the contents of available memory

- core attention directly present
- pre-conscious attention possible to attend to next cycle
- unconscious present in memory and available only later

In the TM the core (cell where the head is) is also available in the next cycle

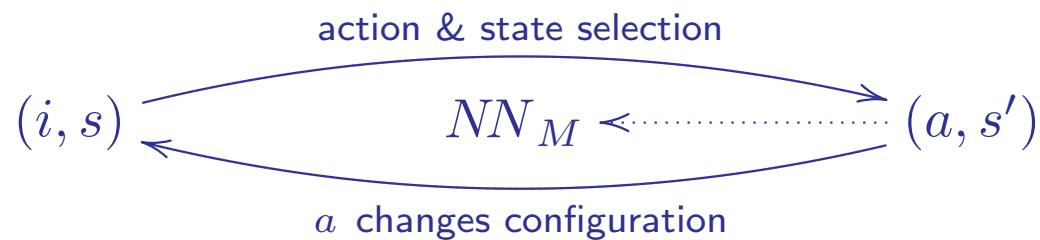
In human cognition better consider core and pre-conscious attention as overlapping

Our conscious attention can be pulled away by a state with desire or aversion

The NN brain has evolved via genes

The TM brain runs with memes, often hampered by states (emotions)

Hybrid Turing model for human cognition [1], [2]



Colleague molecular biology: *“Model also applies to molecular mechanisms: discreteness and states (switching on/off genes) ‘core-attention’ (produced proteins for direct use) ‘pre-attention’ (prepared proteins for later use) ‘un-attention’ (potential proteins dormant in genome)”*

Discreteness (timescales in ms 10 – 30, 100 – 300, 1000 – 3000 ms)

attentional blink	(10-30)
psychological refractory period	(10-30)
memory search retrieval	(10-30)
wagon-wheel illusion	(100-300)
phenomenology	(1000-3000)

Forms of attention

core-conscious

pre-conscious

un-conscious ('subliminal conscious')

Operational definitions given in [0]

States

Mathematical necessity

Overcoming biological noise efficiently [1]

Crucial role for states (including emotions) ( $i, s$ )

Need for state-change and state-preservation:

many neuropeptides ( $> 50$ )

volume transmission (CSF [3], vasopressin, oxytocin [4],  $\beta$ -endorphin [5])

Fundamental instability of perceptible objects (*dukkha*)

Description of meditation

unwholesome states depending on stabilizing peptides (addiction)

goal decreasing frequency of unwholesome states

method sensory restriction

attentional focus on instantaneous configuration

concentration (making core- and preconscious attention coincide)

reflection (taking distance from emotions and other states)

observing stream of states

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[0] Dehaene, Changeux, Naccache, Sackur, Sergent  
Conscious, preconscious, and subliminal processing: a testable taxonomy  
*Trends in Cognitive Sciences*, 10(5), 2006, 204-211.

[1] Zylberberg, Dehaene, Roelfsema, Sigman  
The human Turing machine: a neural framework for mental programs  
*Trends in Cognitive Sciences*, 2011, 15(7), 293-300.

[2] Barendregt, Raffone  
Conscious cognition as a discrete, deterministic, and universal Turing Machine process  
Alan Turing - His Work and Impact. Eds. Cooper and van Leeuwen, Elsevier, 2012.

[3] Veening, Barendregt  
The regulation of brain states by neuroactive substances distributed via the cerebrospinal fluid. A review *Cerebrospinal Fluid Research*. 7(1), 2010.

[4] Veening, de Jong, Barendregt  
Oxytocin messages via the cerebrospinal fluid: behavioral effects; a review  
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[5] Veening, Gerrits, Barendregt (In preparation)  
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